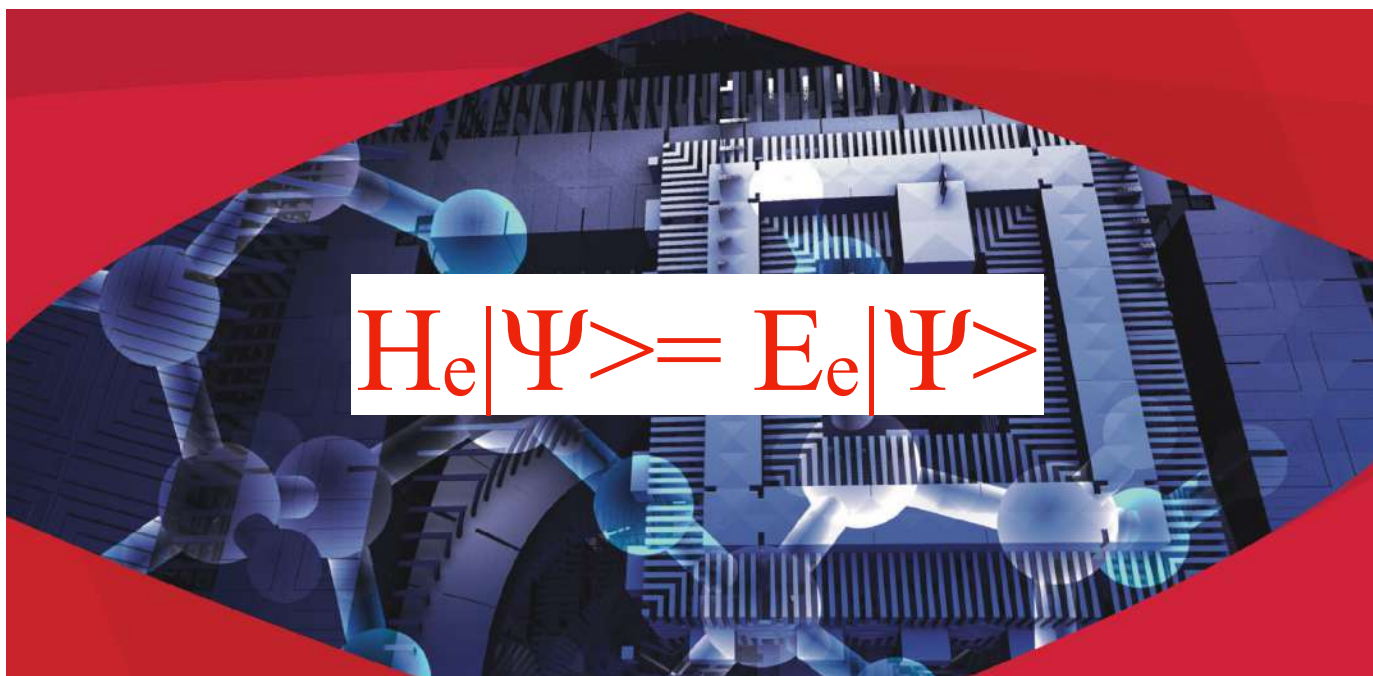


# Quantum Chemistry on a Quantum Computer

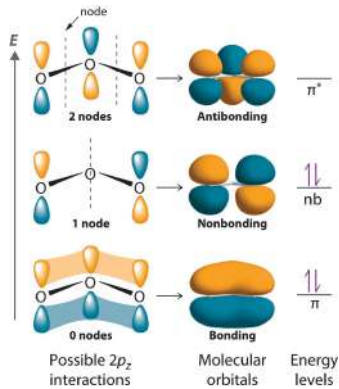


Artur Izmaylov



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# New hope: Digital quantum computers

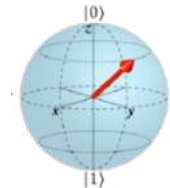


$$\hat{H}_e(\mathbf{R}) |\Phi_j(\mathbf{R})\rangle = E_j(\mathbf{R}) |\Phi_j(\mathbf{R})\rangle$$

$$|\Phi_j(\mathbf{R})\rangle = \sum_{\bar{n}} C_{\bar{n},j}^{2^{N_o}} |n_1, n_2, \dots, n_{N_o}\rangle, \quad n_i = \{0, 1\}$$

For  $|\Psi\rangle$ : occupations to qubit orientations

$$|1100\rangle \rightarrow |\uparrow\uparrow\downarrow\downarrow\rangle$$



## Entanglement

$$|\Psi\rangle = \alpha|\uparrow\uparrow\rangle + \beta|\downarrow\uparrow\rangle + \gamma|\uparrow\downarrow\rangle + \delta|\downarrow\downarrow\rangle$$

N qubits can store  $2^N$  complex numbers



# Quantum Algorithms

“Nature isn’t classical, dammit, and if you want to make a simulation of nature, you'd better make it quantum mechanical...”

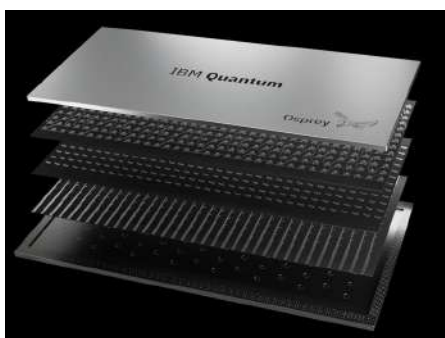
$$|\text{answer}\rangle = U|\text{question}\rangle$$

$$U = U_1 \dots U_n$$

...and by golly it's a wonderful problem, because it doesn't look so easy.”

# Quantum Computers: What can they do?

Short answer: **not much** (in spite of all the hype)

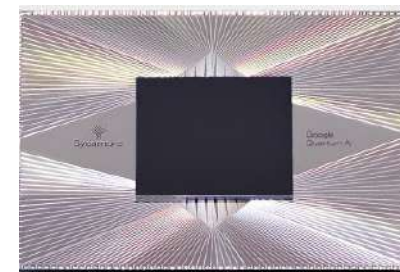


IBM 433Q

$$|\text{answer}\rangle = U|\text{question}\rangle$$

$$U = U_1 \dots U_n$$

**Error correction is needed**



Google 72Q

Error correction: 1-10K qubits for 1 “corrected qubit” ~ **1M qubits**

Expectations: hardware will come, develop algorithms for quantum advantage

Another possibility: quantum inspired methods for classical computers

# Quantum Advantage in Quantum Chemistry

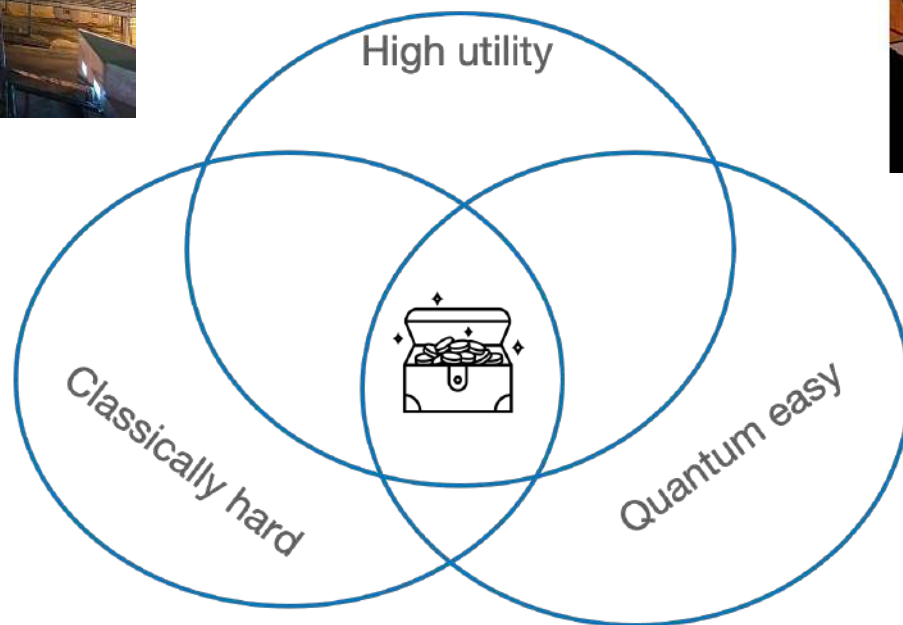


Classical heuristics:  
inaccurate (DFT) or  
expensive (DMRG,  
MRCI, stochastic  
methods,...)

Chemical space:

- Catalysts
- OLEDs

High utility



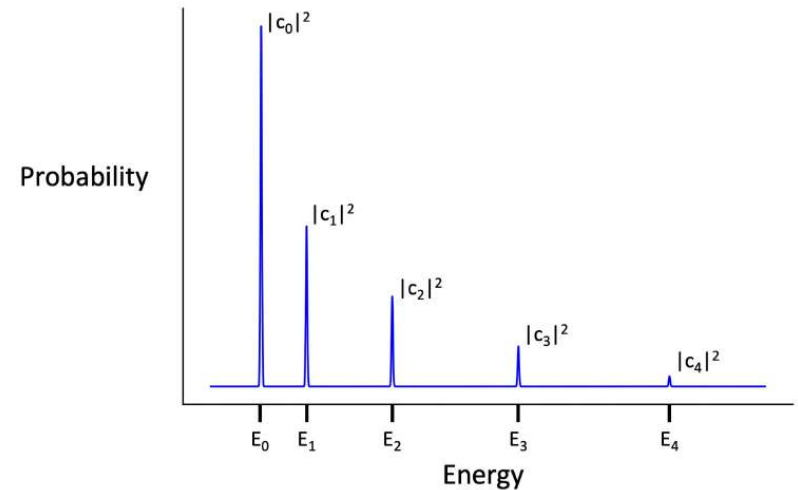
Quantum methods:

- What makes system “quantum easy”?
- Different algorithms different complexities

# Quantum Phase Estimation

$$\hat{H} |\phi_n\rangle = E_n |\phi_n\rangle$$

$$|\psi\rangle = \sum_n c_n |\phi_n\rangle$$



$$\int \langle \psi | e^{-i\hat{H}t} | \psi \rangle e^{i\omega t} dt = \sum_n |c_n|^2 \delta(E_n - \omega)$$

Q1: How to prepare the initial state?

Q2: How to propagate?

# Outline

$$\int \langle \psi | e^{-i\hat{H}t} | \psi \rangle e^{i\omega t} dt = \sum_n |c_n|^2 \delta(E_n - \omega)$$

- Variational Quantum Eigensolver (near-term / quantum inspired)
  - State preparation: Qubit Coupled Cluster ansatz
  - Measurement
- Quantum Phase Estimation (assuming error correction)
  - Trotterization (cost  $\sim t^2$ )
  - Linear Combination of Unitaries (cost  $\sim t$ )

# Fermion-qubit mapping

$$\hat{H}_e = \sum_{pq} h_{pq} \hat{a}_p^\dagger \hat{a}_q + \sum_{pqrs} g_{pq,rs} \hat{a}_p^\dagger \hat{a}_q^\dagger \hat{a}_r \hat{a}_s$$

Polynomial scaling of # of integrals with system size is another reason to use quantum computing

Jordan-Wigner (1928):

$$\hat{a}_j \rightarrow (\hat{x}_j - i\hat{y}_j) \otimes \hat{z}_{j-1} \dots \otimes \hat{z}_1$$

$$\hat{a}_j^\dagger \rightarrow (\hat{x}_j + i\hat{y}_j) \otimes \hat{z}_{j-1} \dots \otimes \hat{z}_1$$

$$\hat{x} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

$$\hat{y} = \begin{pmatrix} 0 & i \\ -i & 0 \end{pmatrix}$$

$$\hat{z} = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

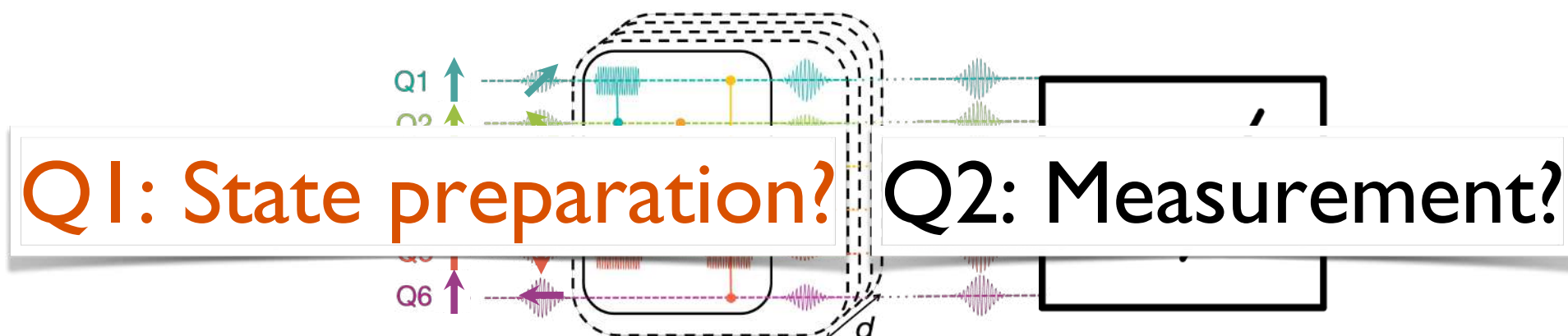
Iso-spectral to  $H_e$  qubit-Hamiltonian:

$$\hat{H}_q = \sum_k c_k \hat{P}_k, \quad \hat{P}_k = \prod_i \hat{\sigma}_i, \quad \hat{\sigma}_i = \{\hat{x}_i, \hat{y}_i, \hat{z}_i\}$$



# I. Variational Quantum Eigensolver

$$E_e = \min_{\theta} \langle 0 | U(\theta)^\dagger H_q U(\theta) | 0 \rangle$$



A. Anand *et al. Chem. Soc. Rev.* **51**, 1659 (2022)

T.C. Yen *et al. npj Quant. Inf.* **9**, 14 (2023)

$|0\rangle$

$U(\theta)|0\rangle$

$H_q$   
measurements

classical  
computer

$$\min_{\theta} \langle 0 | U(\theta)^\dagger H_q U(\theta) | 0 \rangle$$

Alan Aspuru-Guzik *et al. Nat. Commun.* **5**, 4213 (2014)

More on VQE: *Chem. Rev.* **2019**, 119, 10856

*Rev. Mod. Phys.* **2020**, 92, 015003

# I.I State preparation

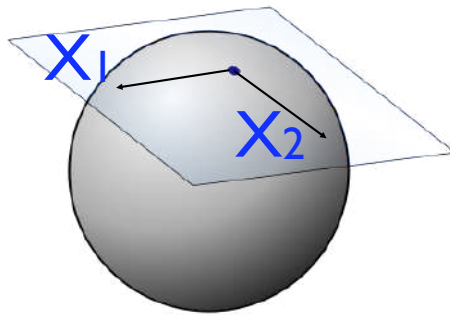
$U_e|\Phi_e\rangle$  (fermions)  $\rightarrow$   $U_q|\Phi_q\rangle$  (qubits)  $\rightarrow$  Sequence of gates  $|0\rangle$   
 (hardware-efficient)

$$\hat{U}(\bar{\theta}) = \prod_k e^{\theta_k \hat{X}_k}, \quad \hat{X}_k^\dagger = -\hat{X}_k$$

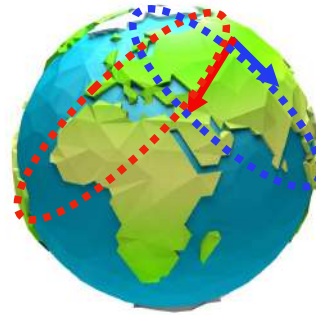
Which  $X$ 's and in what order?

$H = H_0 + V:$

$l$



$\hat{I}_j =$



$\hat{A}_+$

$[\hat{H}_0, \hat{A}_\pm] = \lambda_\pm \hat{A}_\pm$   
 raising/lowering

Energy gradient wrt  $X$ 's amplitude ( $\theta$ ):  $\frac{d}{d\theta} \langle \Phi | e^{-\theta \hat{X}} \hat{H} e^{\theta \hat{X}} | \Phi \rangle \Big|_{\theta=0} = \langle \Phi | [\hat{H}, \hat{X}] | \Phi \rangle$

# Fermionic and Qubit Algebras

$$\hat{U}(\bar{\theta}) = \prod_k e^{\theta_k \hat{X}_k}, \quad \hat{X}_k^\dagger = -\hat{X}_k$$

## Fermionic algebra

$|\Phi_e\rangle$  Hartree-Fock

$$\hat{X}_k = \hat{a}_a^\dagger \hat{a}_i - \hat{a}_i^\dagger \hat{a}_a, \hat{a}_a^\dagger \hat{a}_b^\dagger \hat{a}_i \hat{a}_j - \hat{a}_j^\dagger \hat{a}_i^\dagger \hat{a}_b \hat{a}_a, \dots$$

Methods: UCC, ADAPT-VQE, GUCC...



symmetry conservation  
gate cost

A. Anand et al. *Chem. Soc. Rev.* 51, 1659 (2022)

## Qubit algebra

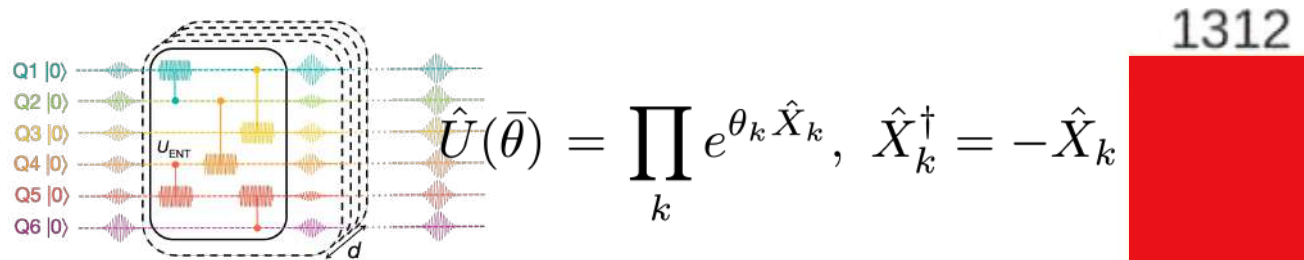
$|\Phi_q\rangle$  simple product state

$$\hat{X}_k = i \bigotimes_{j=1}^N \hat{\sigma}_j, \quad \hat{\sigma}_j = \hat{x}_j, \hat{y}_j, \hat{z}_j, \hat{1}_j$$

Methods: QCC, Qubit-ADAPT, ...  
Ryabinkin et al. *JCTC* 16, 1055 (2020)

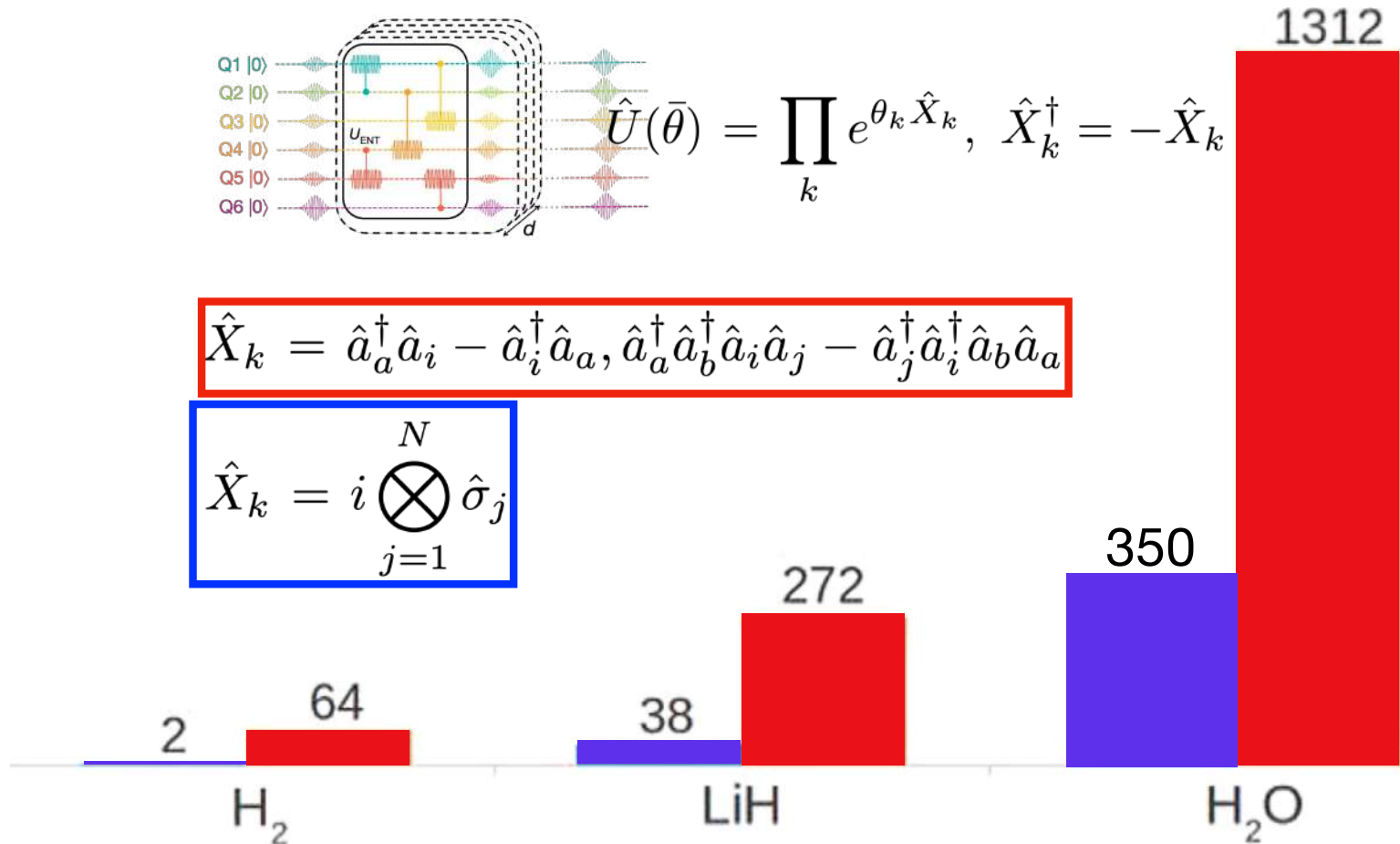


# Two-qubit gate count: QCC vs UCCSD



$$\hat{X}_k = \hat{a}_a^\dagger \hat{a}_i - \hat{a}_i^\dagger \hat{a}_a, \hat{a}_a^\dagger \hat{a}_b^\dagger \hat{a}_i \hat{a}_j - \hat{a}_j^\dagger \hat{a}_i^\dagger \hat{a}_b \hat{a}_a$$

$$\hat{X}_k = i \bigotimes_{j=1}^N \hat{\sigma}_j$$



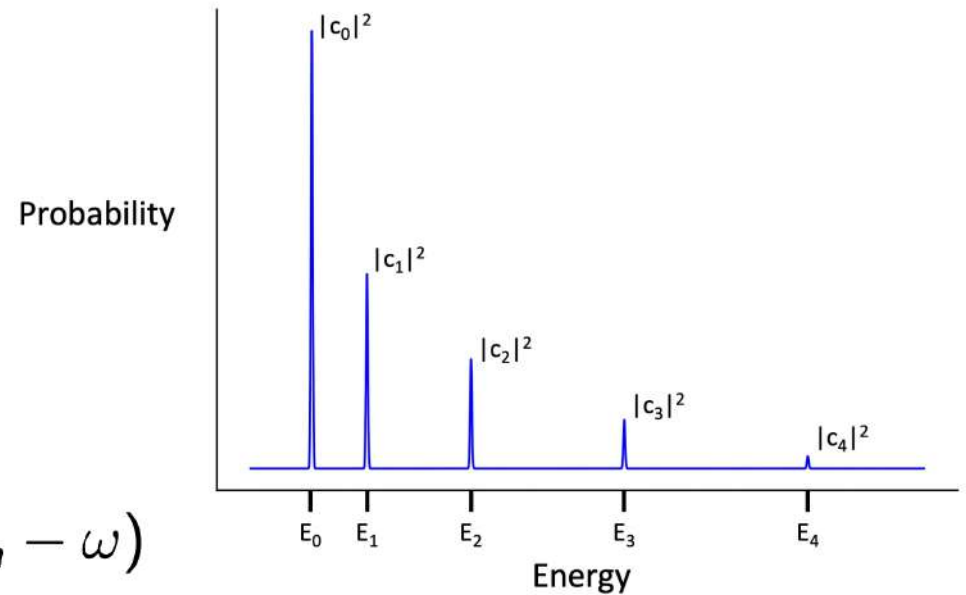
Ryabinkin et al. *JCTC* 14, 6317 (2018)

## 2. Unitary evolution

$$\hat{H} |\phi_n\rangle = E_n |\phi_n\rangle$$

$$|\psi\rangle = \sum_n c_n |\phi_n\rangle$$

$$\int \langle \psi | \underbrace{e^{-i\hat{H}t}}_{\text{blue}} | \psi \rangle \underbrace{e^{i\omega t}}_{\text{red}} dt = \sum_n |c_n|^2 \delta(E_n - \omega)$$



**Unitaries!** Can be done on a quantum computer

## 2.1 Trotterization

$$\boxed{e^{-it\hat{H}}} = e^{-it \sum_n \hat{H}_n} = \lim_{K \rightarrow \infty} \left( \prod_n e^{-i \frac{t}{K} \hat{H}_n} \right)^K$$

$\hat{H}_n$  Hamiltonians: easy to bring to the diagonal form

$$\hat{H}_n = \hat{U}_n^\dagger \hat{Z}_n \hat{U}_n \quad e^{-it\hat{H}_n} = \hat{U}_n^\dagger e^{-it\hat{Z}_n} \hat{U}_n$$

Diagonal:

$$\hat{Z}_n = \sum_i a_i \hat{z}_i + \sum_{ij} b_{ij} \hat{z}_i \hat{z}_j + \dots \quad \text{qubits}$$

$$\hat{Z}_n = \sum_i c_i \hat{n}_i + \sum_{ij} d_{ij} \hat{n}_i \hat{n}_j + \dots \quad \text{fermions}$$

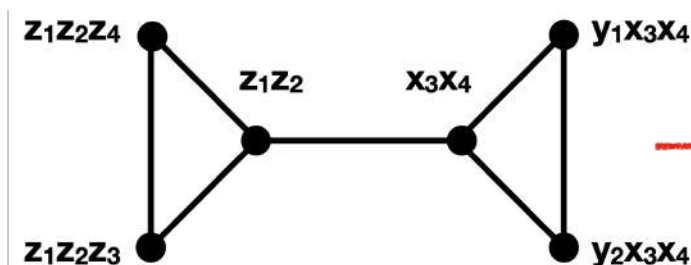
$\hat{n}_i = \hat{a}_i^\dagger \hat{a}_i$

# What Hamiltonians can we rotate to the diagonal form?

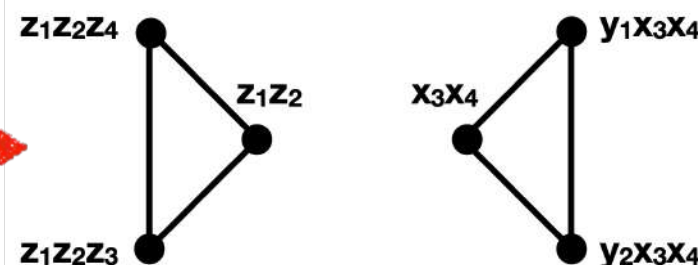
## Qubit algebra: Commuting Pauli products

(Qubitization-then-partitioning)

$$\hat{H}_n = \sum_k c_k \hat{P}_k, \quad \hat{H} = \sum_n \hat{H}_n$$
$$[\hat{P}_k, \hat{P}_{k'}] = 0 : \quad \hat{U}_{\text{Cliff}}^\dagger \hat{H}_n \hat{U}_{\text{Cliff}} = \hat{Z}_n$$



Verteletskyi, et al. *J. Chem. Phys.* 152 124114 (2020)



Yen, et al. *J. Chem. Theory Comput.* 16, 190 (2020)

# Hamiltonian decomposition: Fermionic algebra

(Partitioning-then-qubitization)

$$\hat{H} = \sum_n \hat{H}_n \quad \text{Hartree-Fock solvable fragments}$$

**1e:**  $\hat{h}_{1e} = \sum_{pq} h_{pq} \hat{a}_p^\dagger \hat{a}_q = \hat{U}_{\text{HF}}^\dagger \left( \sum_p \epsilon_p \hat{n}_p \right) \hat{U}_{\text{HF}}$

$\hat{U}_{\text{HF}} = \prod_{p>q} e^{\theta_{pq}(\hat{a}_p^\dagger \hat{a}_q - \hat{a}_q^\dagger \hat{a}_p)}$

occ. number  $\hat{n}_i \rightarrow \hat{z}_i$

orbital rotations

**1e+2e:**  $\hat{H}_{\text{HF}} = \hat{U}_{\text{HF}}^\dagger \left( \sum_{pq} \lambda_{pq} \hat{n}_p \hat{n}_q \right) \hat{U}_{\text{HF}}$



# Which fragments are better, Qubit or Fermionic?

$$\hat{H} = \sum_n \hat{H}_n$$

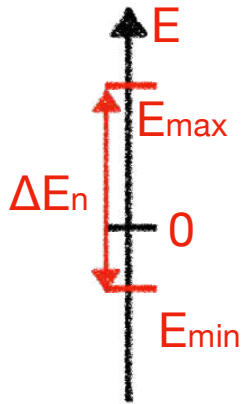
Practical choice: 1st order Trotter

$$e^{-iHt} \approx \prod_n e^{-iH_n t} = \prod_n \hat{U}_n^\dagger e^{-iZ_n t} \hat{U}_n$$

# Optimizing Trotter error

$$\|e^{-i\hat{H}t} - \prod_n e^{-i\hat{H}_n t}\| \leq t^2 \sum_{k \neq n} \|[\hat{H}_k, \hat{H}_n]\|$$

A.Childs et al. PRX 11, 011020 (2021)



$$2 \sum_{n \neq k} \|[\hat{H}_n, \hat{H}_k]\| \leq \left( \sum_n \Delta E_n \right)^2 \left[ 1 - \sum_n \omega_n^2 \right]$$

Total Spectral  
Range

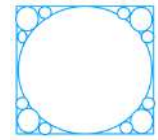
Entropic  
part

$$S = - \sum_n \omega_n \ln(\omega_n) \approx 1 - \sum_n \omega_n^2 = S_L$$

$$\omega_n = \frac{\Delta E_n}{\sum_k \Delta E_k} : \sum_n \omega_n = 1, \omega_n \in [0, 1]$$

Smaller spread of  
eigenvalues

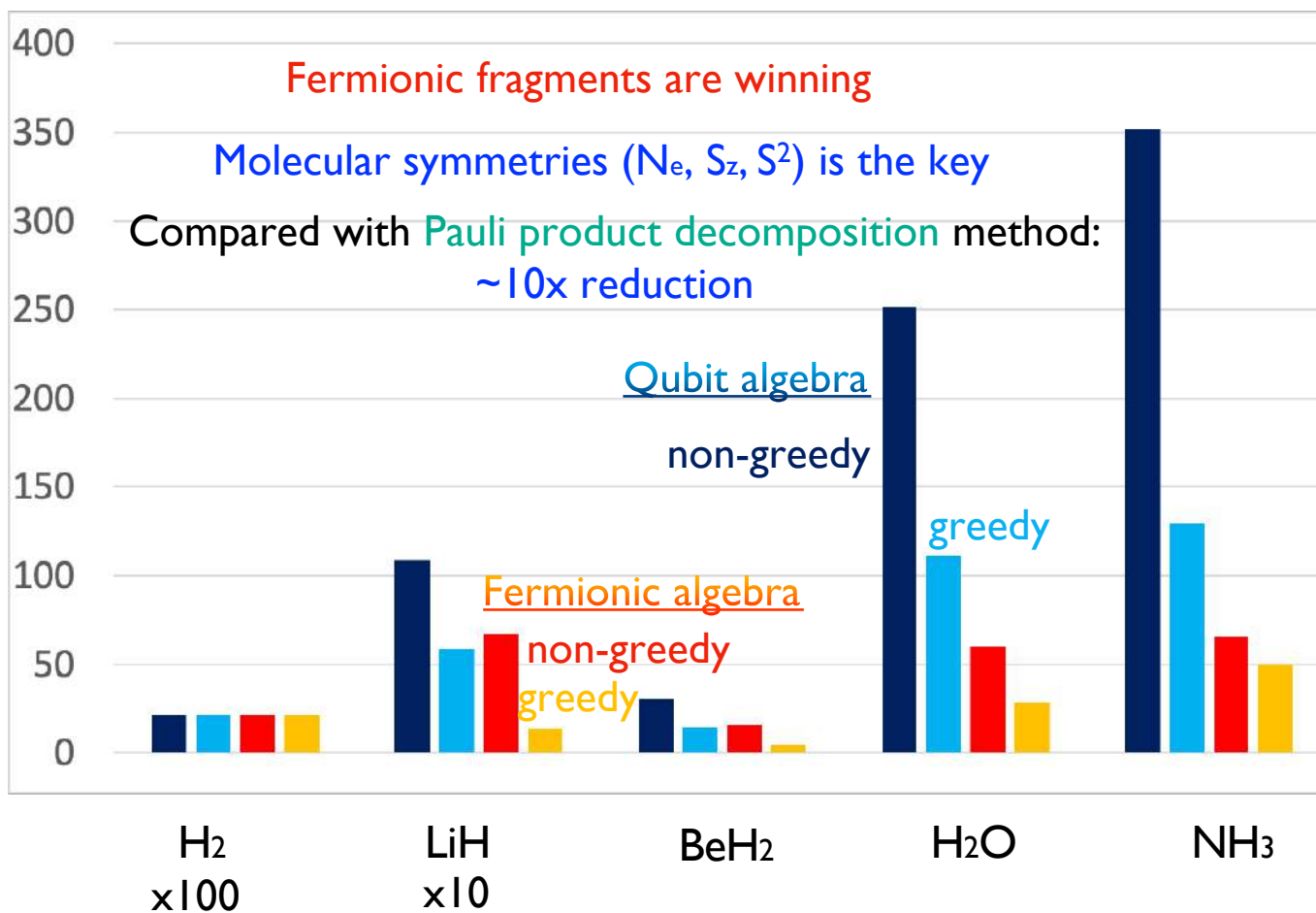
Non-uniform  
distribution



“Greedy” algorithms

Quantum 7, 1086 (2023)

# Trotterization error



$$\sum_{nn'} ||[H_n, H_{n'}]||$$

## 2.2 Linear Combination of Unitaries

$$\hat{H} = \sum_k c_k \hat{U}_k$$

$$e^{-i\hat{H}t} = \mathbb{1} - i\hat{H}t + \dots$$

### 1. Taylor expansion

PRL 114, 090502 (2015) PHYSICAL REVIEW LETTERS week ending 6 MARCH 2015

#### Simulating Hamiltonian Dynamics with a Truncated Taylor Series

Dominic W. Berry,<sup>1</sup> Andrew M. Childs,<sup>2,3,4,5</sup> Richard Cleve,<sup>2,5,6</sup> Robin Kothari,<sup>2,6,7</sup> and Rolando D. Somma<sup>8</sup>

### 2. Qubitization

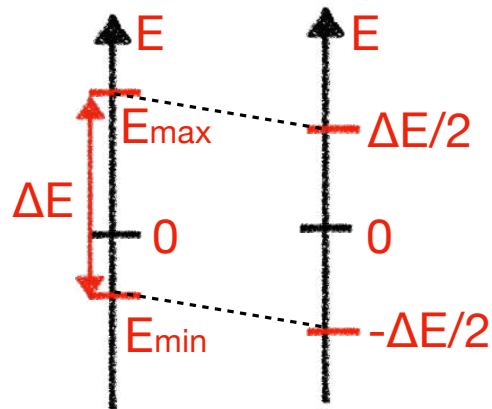
Hamiltonian Simulation by Qubitization

Guang Hao Low<sup>1</sup> and Isaac L. Chuang<sup>2</sup>

$$\hat{W} = \begin{pmatrix} \hat{H} & \sqrt{1 - \hat{H}^2} \\ -\sqrt{1 - \hat{H}^2} & \hat{H} \end{pmatrix} = e^{i \cos^{-1}(\hat{H}) \hat{\sigma}_y}$$

LCU game: reduce the 1-norm of coefficients,  $\sum_k |c_k| \sim \text{cost}$

# I-norm lower bound



$$\hat{H} = s\hat{I} + \sum_k c_k \hat{U}_k$$

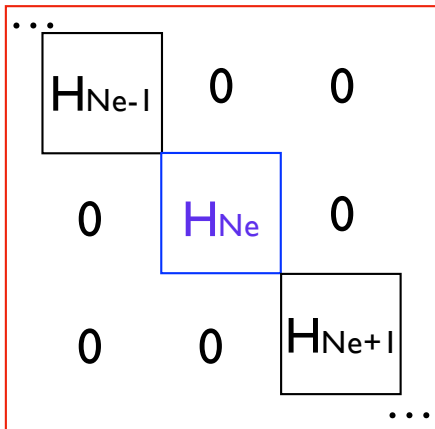
$$\Delta E/2 \leq \|\hat{H} - s\hat{I}\| \leq \sum_k |c_k| \cdot \|\hat{U}_k\| = \sum_k |c_k|$$

$$\frac{E_{\max} - E_{\min}}{2} = \Delta E/2 = \min_s \|\hat{H} - s\hat{I}\|$$

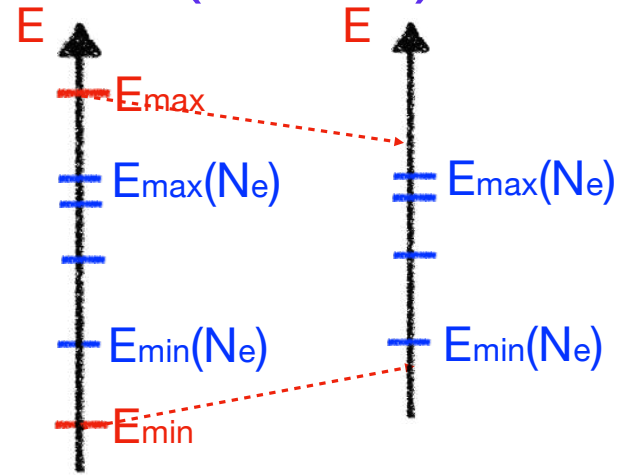
Hard limit: I-norm cannot be lower than a half of the spectral range ( $\Delta E/2$ )

(One way around is to change H)

# Block-Invariant Symmetry Shift (BLISS)



$$\hat{H}\Psi_k = (\hat{H} - \hat{K})\Psi_k = E_k\Psi_k$$



$$\hat{K} = \left[ \alpha_1 + \sum_{pq} \beta_{pq} \hat{a}_p^\dagger \hat{a}_q \right] (\hat{N}_e - N_e) + \alpha_2 (\hat{N}_e^2 - N_e^2)$$

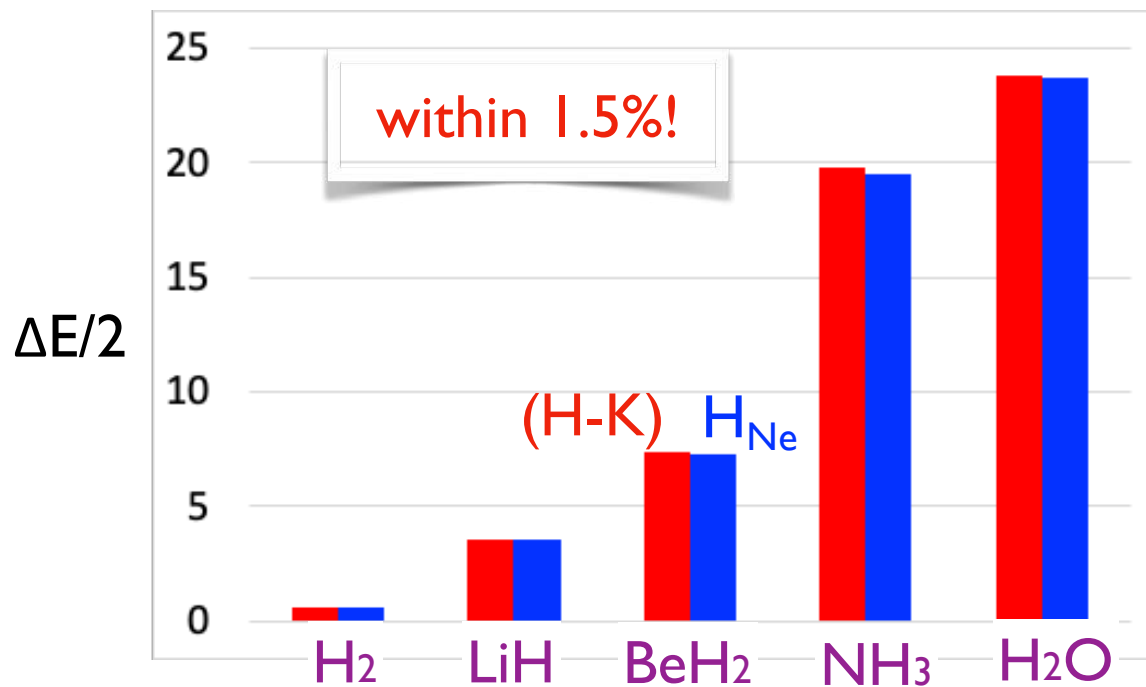
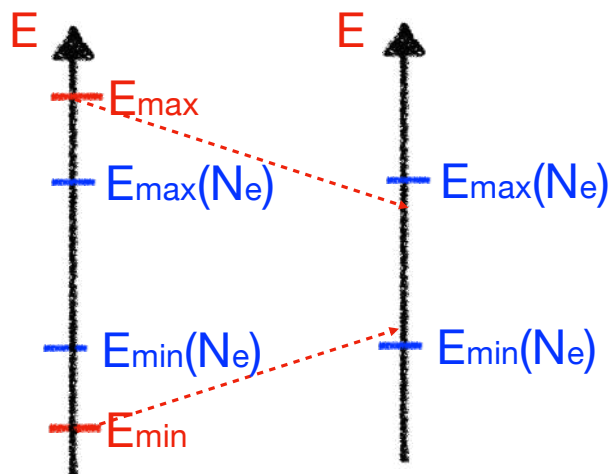
Finding K: optimizing to reduce l-norm

$$(\hat{H} - \hat{K}) = \sum_n c_n \hat{P}_n, \quad \lambda = \sum_n |c_n|$$

Note that 1) K is not a symmetry, 2) using other symmetries is possible

# K search: Lower bound results

Success is when (H-K) has the same  $\Delta E/2$  as in H for a particular # of electrons:



STO-3G basis

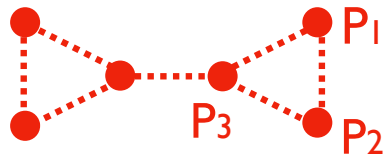
What U's to use?

$$(\hat{H} - \hat{K}) = \sum_n c_n \hat{U}_n$$

P's are good but can we do better?



# Qubit anti-commuting grouping



$$\{\hat{P}_k, \hat{P}_{k'}\} = 2\delta_{kk'}$$

$$\sum_k \tilde{c}_k^2 = 1$$

anti-commuting  
Paulis form  
reflections



$$\hat{R}_i = \sum_k \tilde{c}_k \hat{P}_k$$

Grouping always reduces l-norm

arXiv:2208.08272

## Fermionic reflections:

$$\min_{\hat{U}_{\text{HF}}^{(k)}, \lambda^{(k)}} \|\hat{H} - \sum_k \hat{H}_{\text{HF}}^{(k)}\| \quad \hat{H}_{\text{HF}}^{(k)} = \left(\hat{U}_{\text{HF}}^{(k)}\right)^\dagger \left(\sum_{pq} \lambda_{pq}^{(k)} \hat{n}_p \hat{n}_q\right) \hat{U}_{\text{HF}}^{(k)}$$

$$\hat{n}_p^2 = \hat{n}_p$$

$$\hat{r}_p = 2\hat{n}_p - 1$$

$$\hat{r}_p^2 = 1$$

$$\hat{H} = \sum_k c_k \hat{R}_k$$

1e:  $\hat{R} = \hat{U}_{\text{HF}}^\dagger \hat{r}_p \hat{U}_{\text{HF}}$

2e:  $\hat{R} = \hat{U}_{\text{HF}}^\dagger \hat{r}_p \hat{r}_{p'} \hat{U}_{\text{HF}}$

arXiv:2208.08272

# LCU decompositions results

l-norm improvement w.r.t. naïve Pauli decomposition for H

---

H  $\Delta E/2$  x1.8

H-K  $\Delta E/2$  x3.1

H-K Pauli x1.9

H-K Qubit AC x2.8

H-K Fermionic x2.7

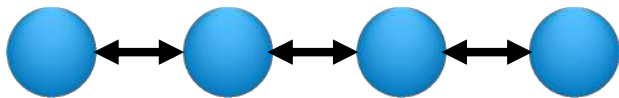
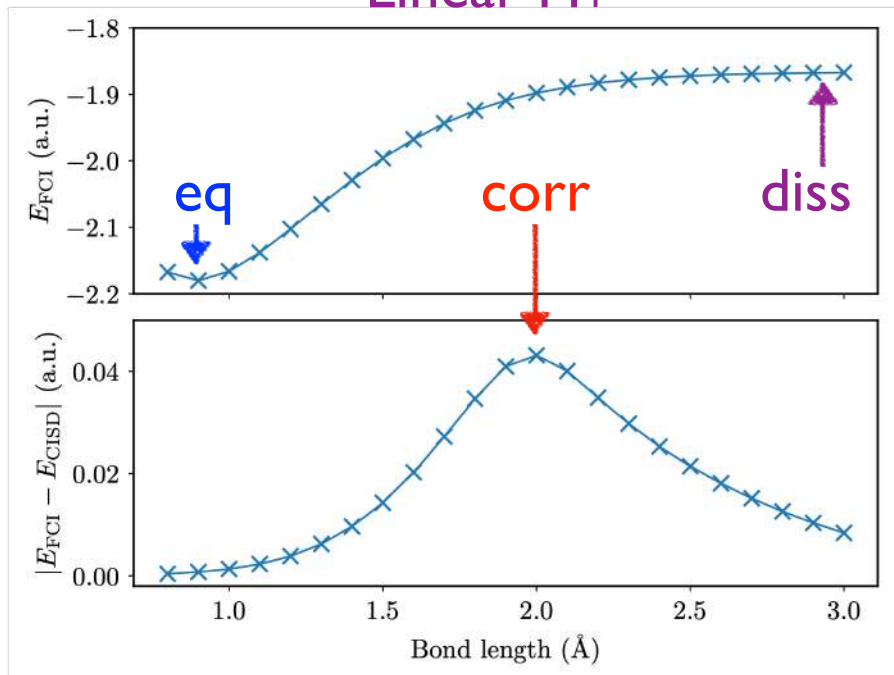
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STO-3G basis H<sub>2</sub> LiH BeH<sub>2</sub> NH<sub>3</sub> H<sub>2</sub>O

arXiv:2304.13772

# Testing complexity: classical vs quantum

## Linear H<sub>4</sub>



## Other systems:

Rectangular H<sub>4</sub>

H<sub>2</sub>O, sym OH stretch

N<sub>2</sub> stretched

## Quantum “hardness”:

1. #  $\exp(\theta_k X_k)$
2. Trotter error
3. LCU cost (1-norm)
4. # measurements VQE

arXiv:2311.00129

# Summary

Quantum chemistry is still attractive for demonstrating quantum advantage: We are looking for classically hard and quantum easy systems

Algebraic techniques are essential for many quantum algorithms due to non-trivial algebraic structure of the Hamiltonian: Lie and Clifford algebras and groups

$$\hat{H} = \sum_n \hat{H}_n \qquad \hat{H} = \sum_k c_k \hat{U}_k$$

Great news is that in all quantum algorithms we have checked so far quantum hardness is not correlated with classical hardness (except initial state preparation)

More info:



Tequila (developed with Alan Aspuru-Guzik group)

# Acknowledgments



Tzu-Ching  
(Thomson) Yen



Aadithya  
Ganeshram



Ignacio Loiza



Robert Lang



Seonghoon  
Choi



Luis Martinez  
Martinez



Zachary  
Pierce Bansingh

Postdoc  
Position  
Opening



Smik  
Patel



Funding:

